DMSN Tutorial 3: Network Centrality and Applications

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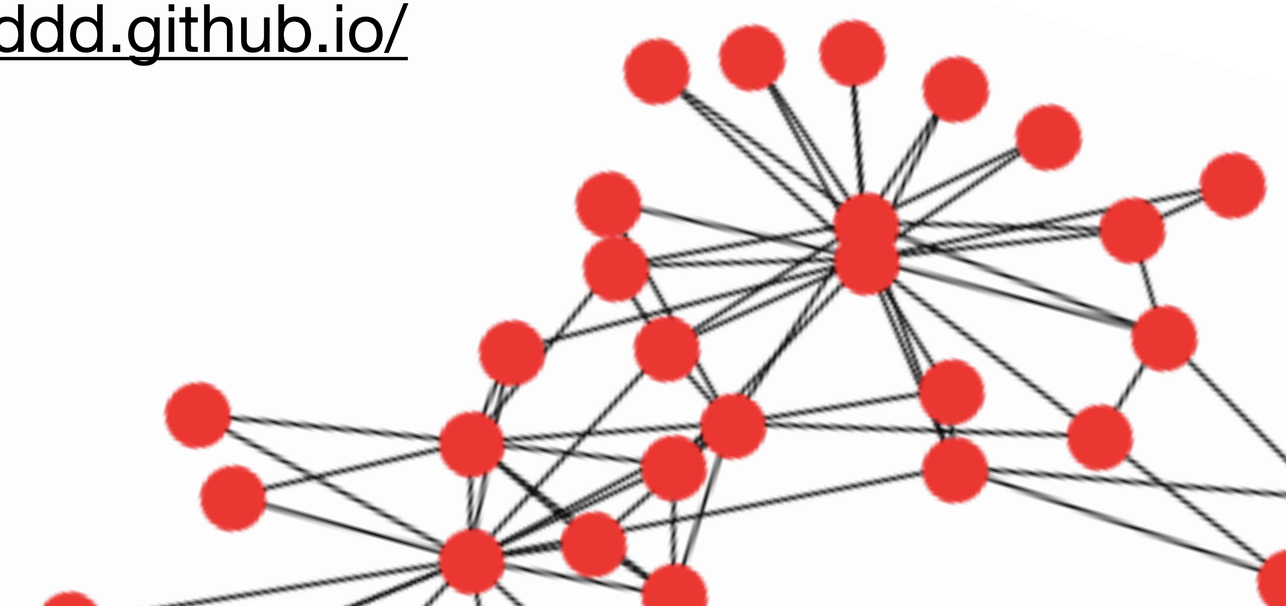
https://narnolddd.github.io/

Morning all! We will start 9:50 to allow for a break after MCQs. Good luck! :)

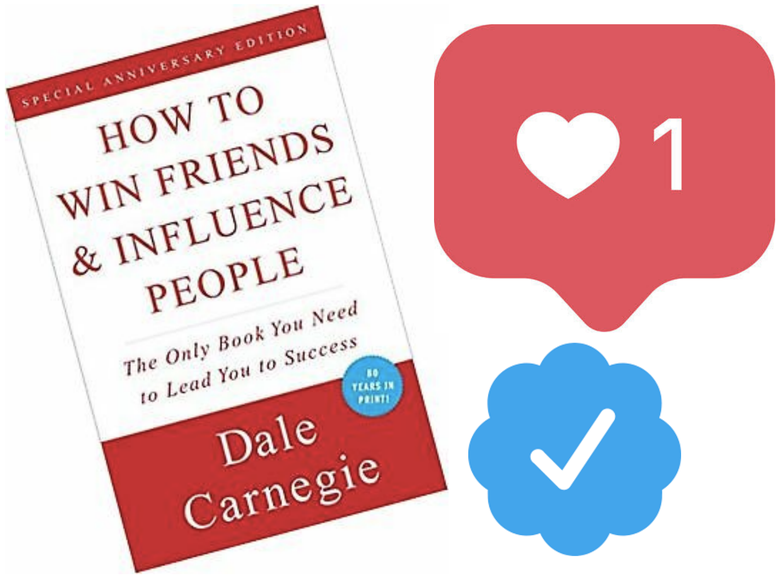


In this tutorial

- **Recap** different centrality measures
- Go over some **centrality calculations**
- Application to the **FIFA 2018** World Cup Final

Who is important in a network?

It depends what you mean by important…

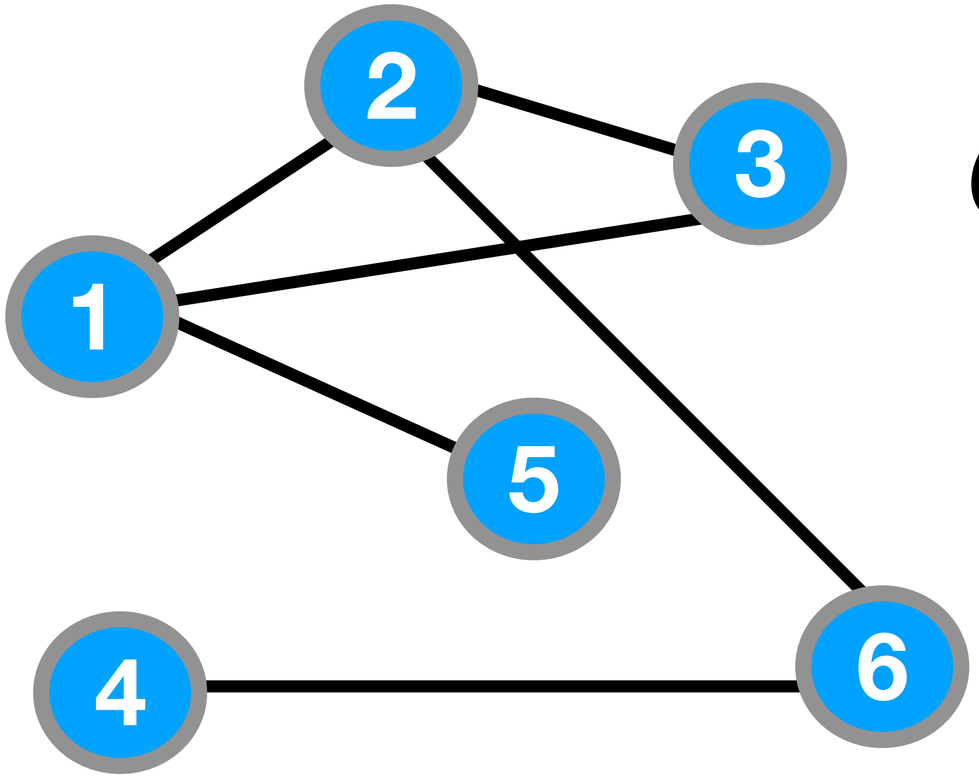


Who has **influence?**
(authority, popularity, …)

Who **facilitates connectivity**/‘glues’ the network together?

(In-)Degree Centrality

Number of connections (followers, friends, citations, …)



$$C(i) = \frac{k_i}{N-1}$$

Degree of node i

Number of **other** nodes in the network, to **normalise** (only really necessary if you're comparing more than one network)

$$C(1) = C(2) = 3/5$$

Freeman Network Centralisation

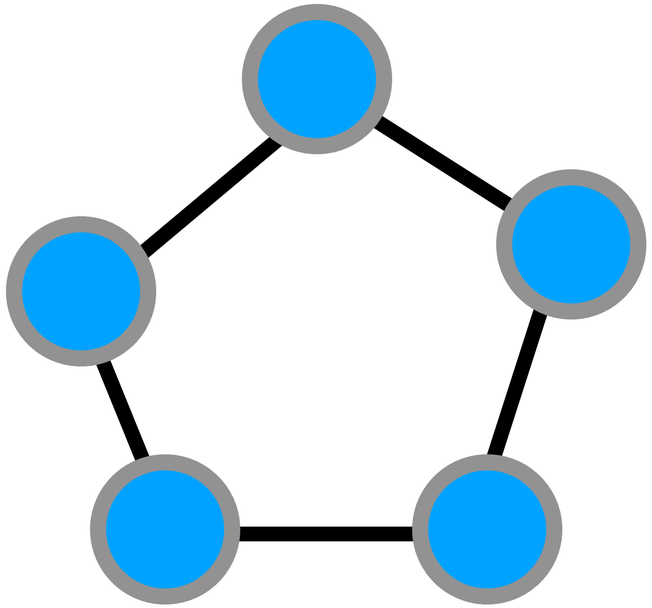
How far from **equal** is a network's degree distribution?

Compare the degree of each node with the largest degree

$$C_F(G) = \frac{\sum_{i=1}^{N} k_{\max} - k_i}{(N-1)(N-2)}$$

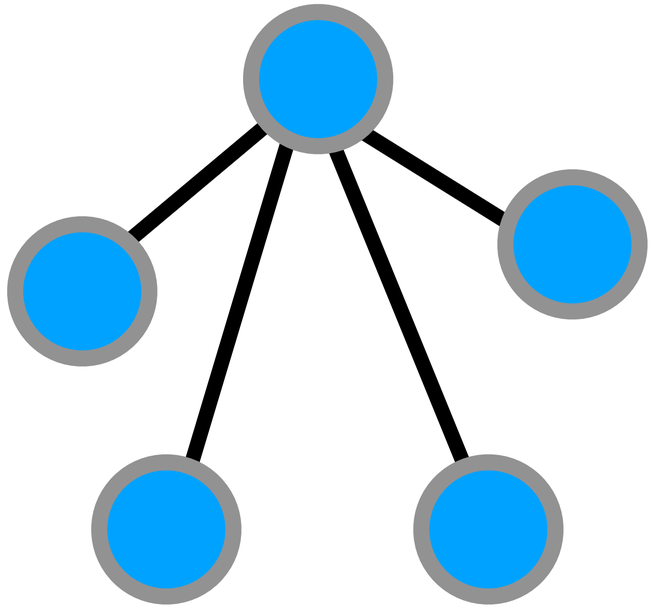
The largest possible value the top could be in a network of **N** nodes

Freeman Network Centralisation



$$C_F = 0$$

Nodes have equal share of the links.



$$C_F = 1$$

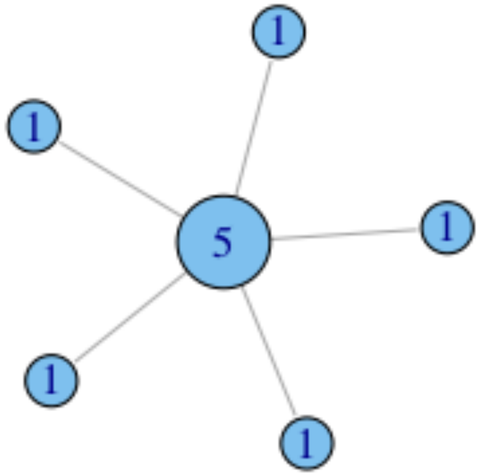
One node has all the links

FREEMAN NETWORK CENTRALITY EXPLAINED

Explanation of the denominator:

In the star topology one node has degree N-1 and all other nodes have degree of 1:

$$0 + ((N-1)-1) * (N-1) \quad = \quad (N-2) * (N-1)$$



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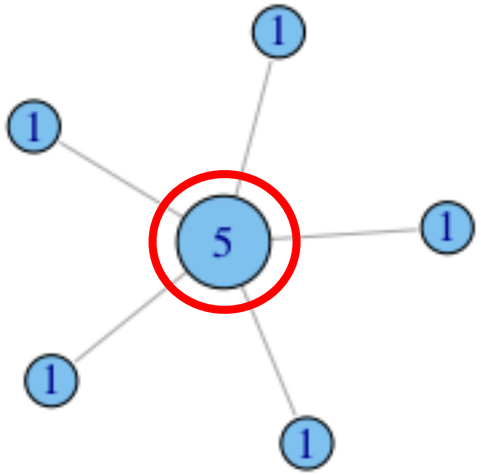

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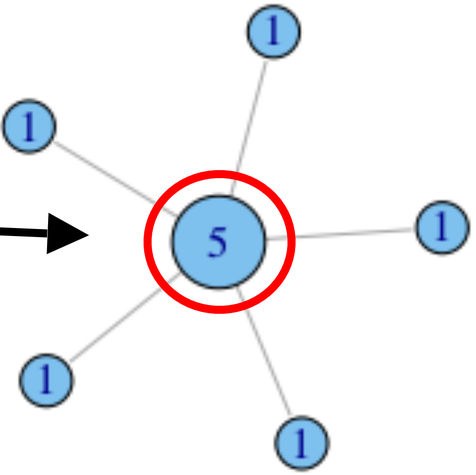

I have N-1 (5) links:
(N-1) - (N-1) = 5 - 5 = 0

FREEMAN NETWORK CENTRALITY EXPLAINED

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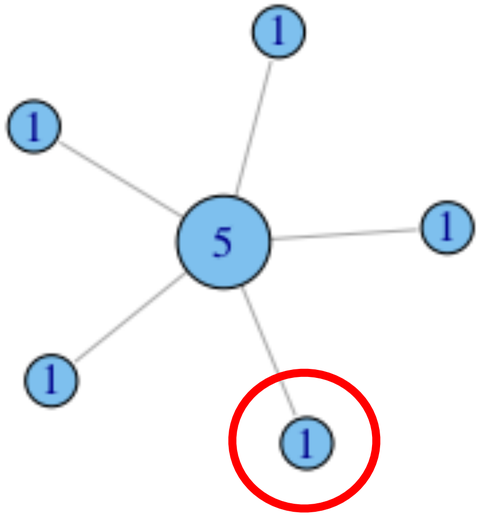
$$0 + ((N-1)-1) * (N-1) \quad = \quad (N-2) * (N-1)$$

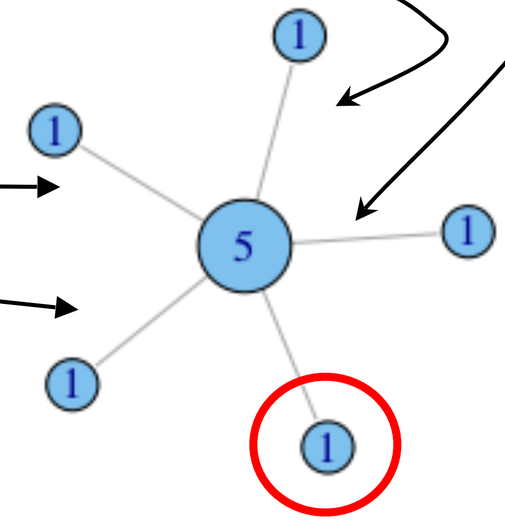


FREEMAN NETWORK CENTRALITY EXPLAINED

Explanation of the denominator:

In the star topology one node has degree N-1 and all other nodes have degree of 1:

$$0 + \mathbf{((N-1)-1)} * (N-1) \quad = \quad (N-2) * (N-1)$$



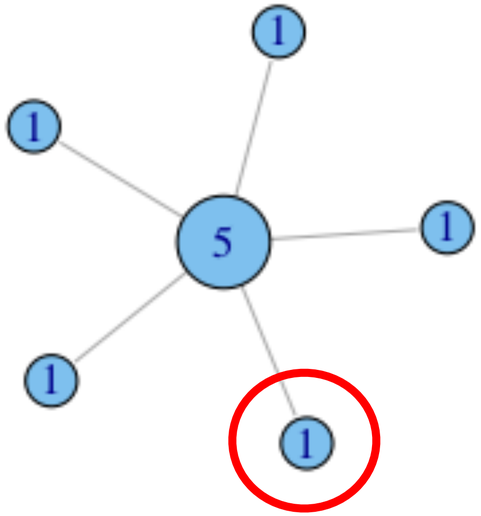
Hub has 4 more links than me

FREEMAN NETWORK CENTRALITY EXPLAINED

Explanation of the denominator:

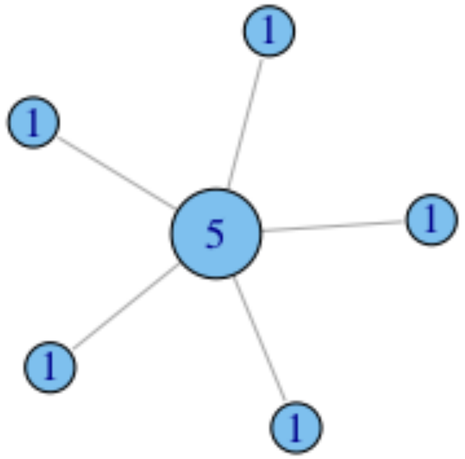
In the star topology one node has degree N-1 and all other nodes have degree of 1:

$$0 + ((N-1)-1) * \mathbf{(N-1)} \quad = \quad (N-2) * (N-1)$$



Need to do the same for all nodes except the hub

FREEMAN'S NETWORK CENTRALITY



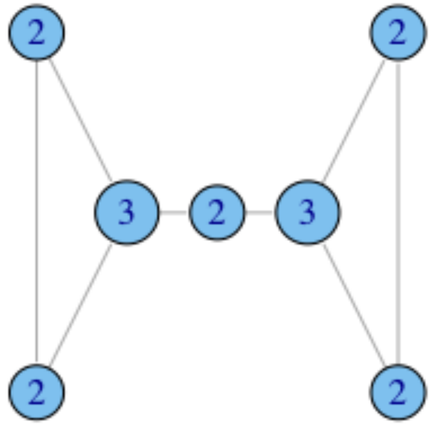
C_D = 1.0

= [4+4+4+4+4]/5*4 = 1

n-1 **n-2**

= [(2-1)+(2-2)+(2-2)+(2-2)+(2-1)]/4*3 = 1/6

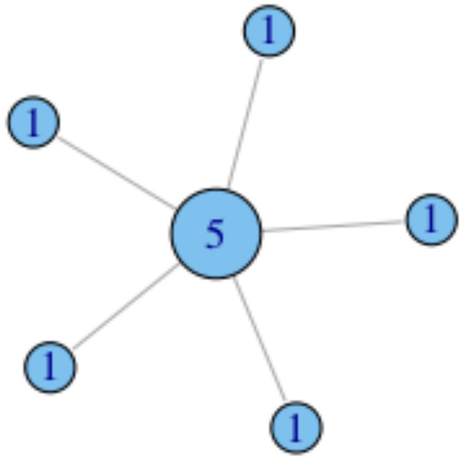
C_D = 0.167



= [(3-2)+(3-2)+(3-3)
+(3-2)+(3-3)+(3-2)
+(3-2)]/(6*5)
= 5/30

C_D = 0.167

FREEMAN'S NETWORK CENTRALITY

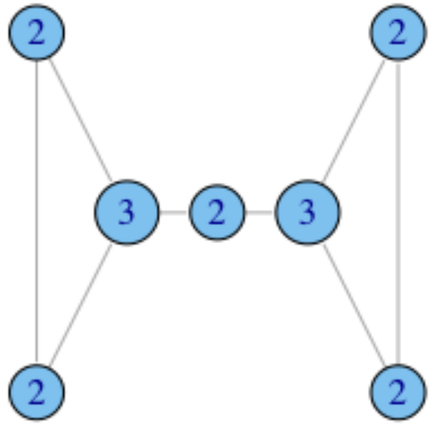


C_D = 1.0

(N=5) 1+0+0+0+1

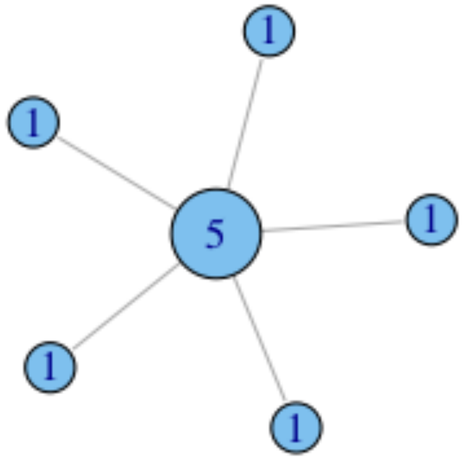
4*3 = 2/12 = 1/6

C_D = 0.167



C_D = 0.167

FREEMAN'S NETWORK CENTRALITY

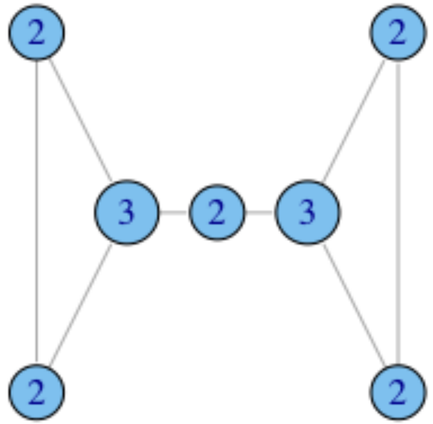


C_D = 1.0

C_D = 0.167

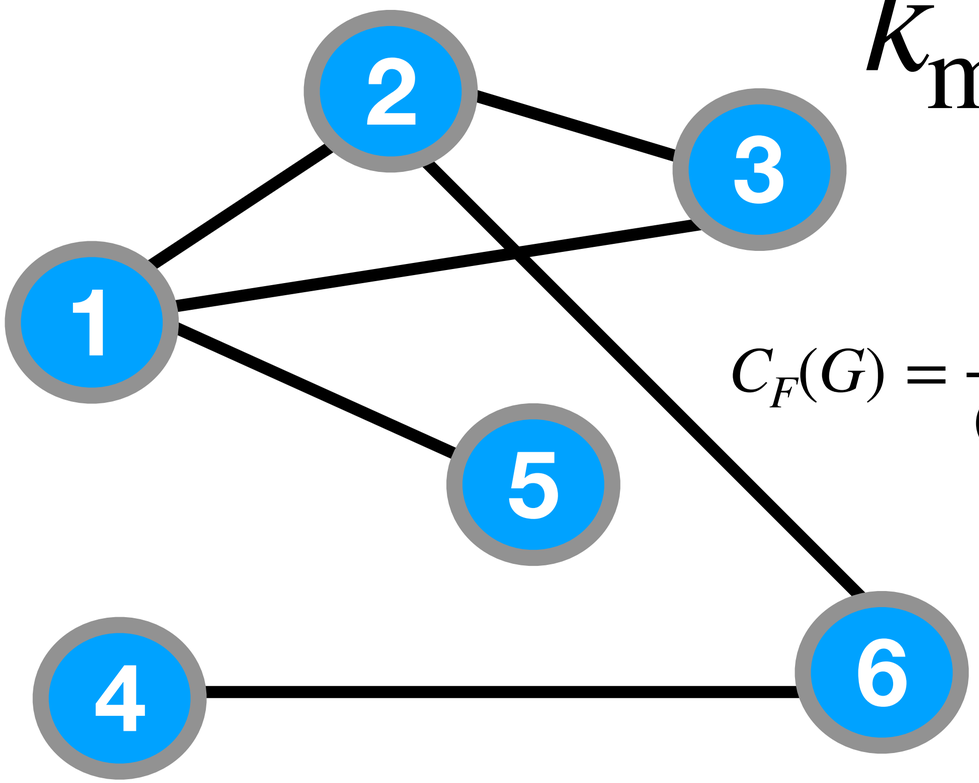
(N=7) 1+1+0+1+0+1+1

6*5 = 5/30



C_D = 0.167

Freeman Network Centralisation



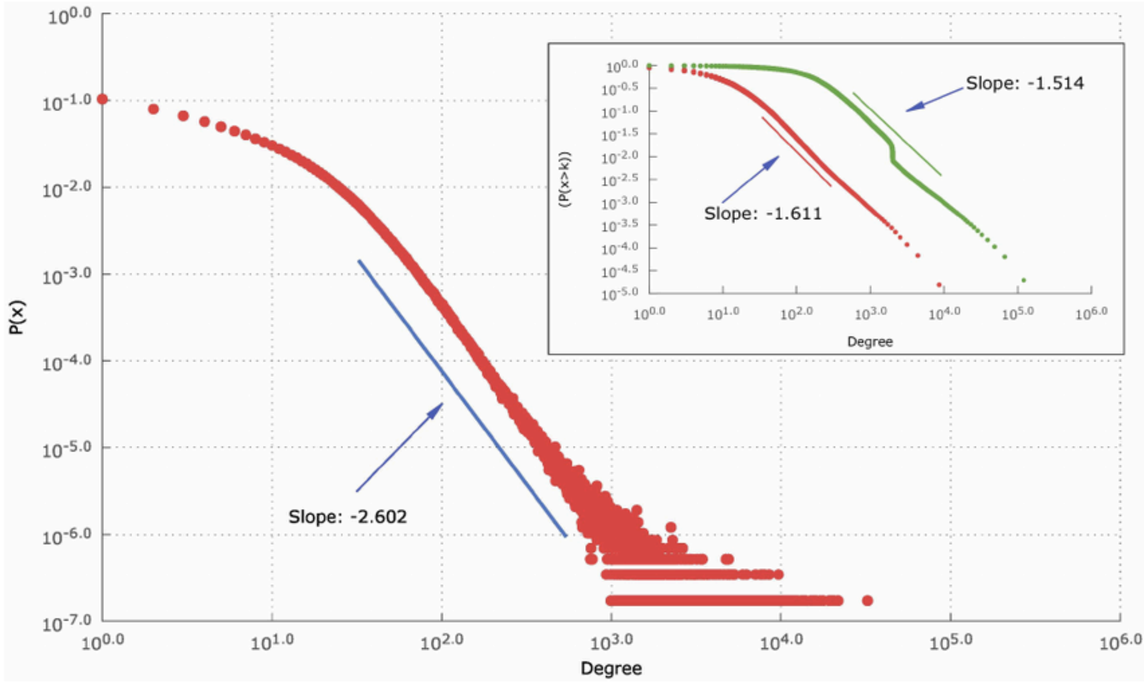
$$k_{\max} = 3$$

$$C_F(G) = \frac{\sum_{i=1}^{N} k_{\max} - k_i}{(N-1)(N-2)}$$

i	k_i	$k_{\max} - k_i$
1	3	0
2	3	0
3	2	1
4	1	2
5	1	2
6	2	1

$$C_F(G) = \frac{0+0+1+2+2+1}{5 \times 4} = \frac{3}{10}$$

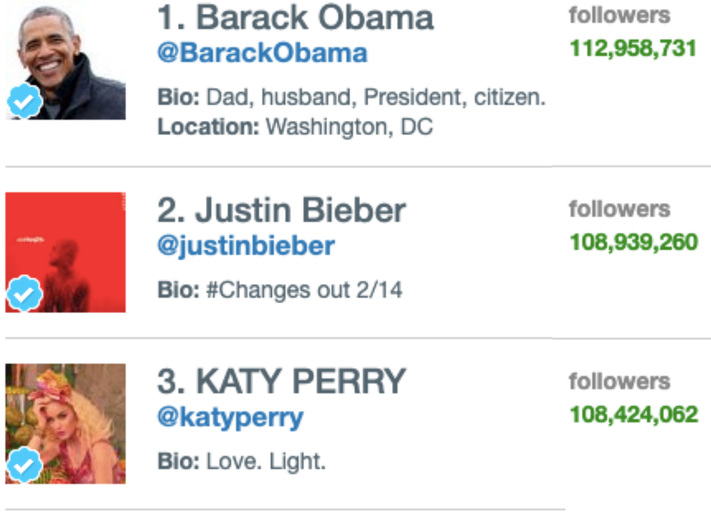
(In-)Degree Centrality: Remarks



Twitter Follower Distribution

[Massive Social Network Analysis: Mining Twitter for Social Good, David Ediger et al (2010)]

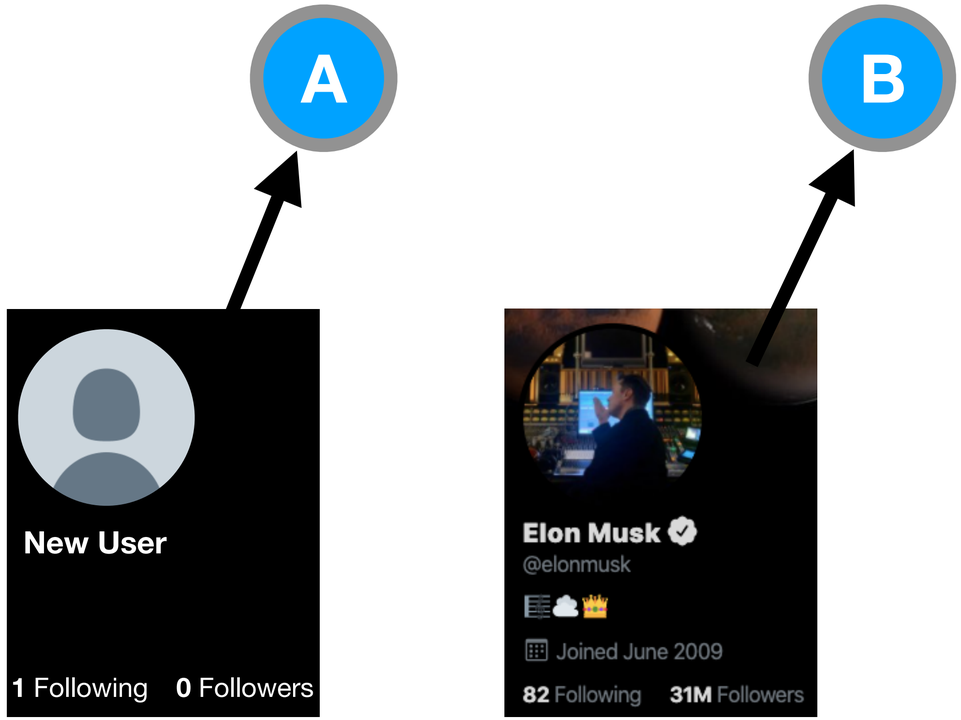
Network (in-)degree distributions are often **heavy tailed**, with a **small number** of nodes having a **huge degree** but most having a **small degree**



Degree centrality **good for ranking top nodes** with extreme values, but **not very good** for ranking nodes in the **middle/low end**

(In-)Degree Centrality: Remarks

Easy metric to **manipulate** in online social networks: buying followers, using bots etc



A and B have **SAME** degree centrality

Gives **every link equal weight** — is this meaningful?

EIGENVECTOR CENTRALITY

1. Assign centrality score of 1 to all nodes
2. **Recompute scores of each node as weighted sum of centralities for a all nodes in a node's neighbourhood**
3. Normalise by dividing each value by the largest value
4. Repeat steps 2 and 3 until centrality values stop changing.

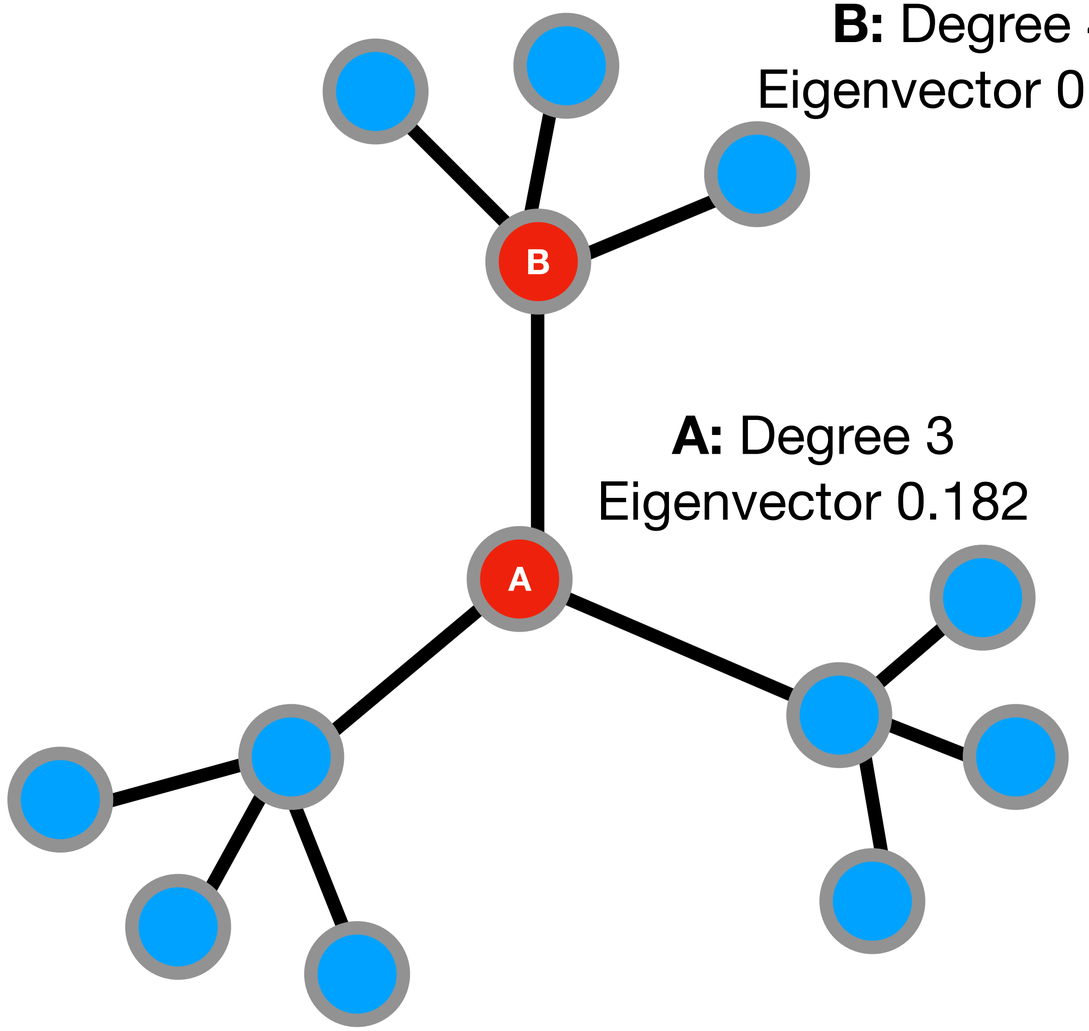
Eigenvector Centrality

Eigenvector centrality addresses **BOTH** of these issues

Each node's centrality is **proportional** to the **sum** of it's neighbours centrality

Node **B** has higher degree centrality **(4 vs 3)**

But **A** has higher eigenvector centrality



Google PageRank

Challenge: how to rank search engine results well, even in the middle?

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Research groups and centres

- **Networks**

networks.eecs.qmul.ac.uk

The Networks Research Group was established in 1987 and is active in key areas of networking including Internet measurements, quality of service, mobile communications, content delivery and network analysis. The group has an international reputation for excellence; our work is regularly published in prestigious venues such as SIGCOMM, INFOCOM, IMC, CoNEXT, WWW, ICNP and various premier IEEE/ACM Transactions (e.g. ToN, TPDS, TC, ToMM). Our research is funded by a mix of grants from EPSRC, EU H2020, and industrial partners. **Current research projects**

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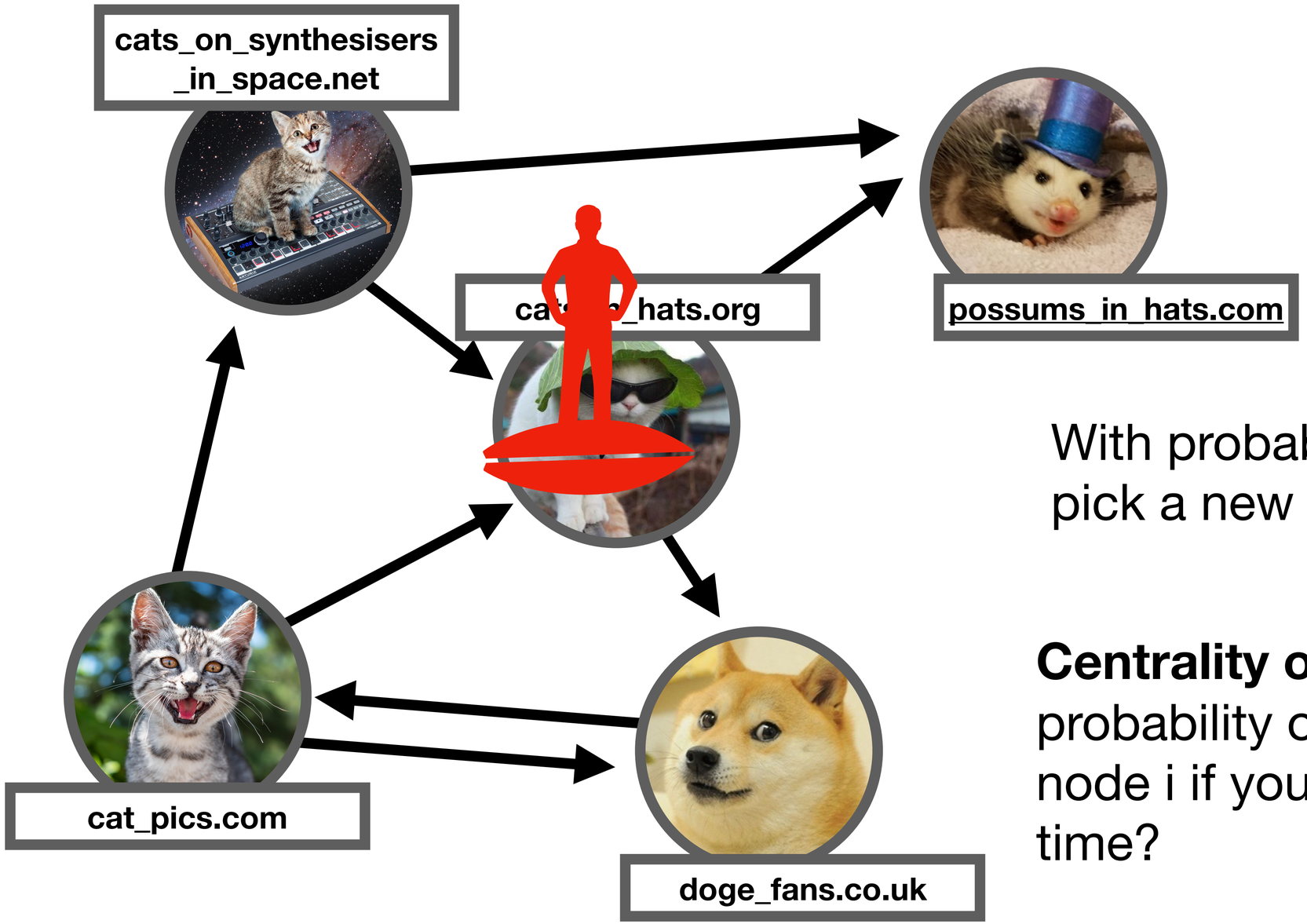
Current Research Activities

- Internet measurements
- Software Defined Networking
- Network programmability
- Topology theory and resilience
- Distributed computing
- Web systems
- Internet of Things
- Digital and social media
- Security

View the Web as a directed graph with web pages as nodes and hyperlinks as links

PageRank idea: the random surfer



"Random surfer" navigates the web by clicking on hyperlinks.

With probability **p** they start over and pick a new webpage to start again

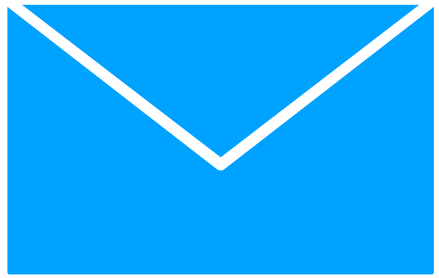
Centrality of node i: what is the probability of finding the surfer at node i if you check after a long time?

Properties of Eigenvector/PageRank

- More difficult metric to **‘cheat’** than degree
- Takes into account full **network structure** and can help **distinguish** nodes in the ‘middle’
- Still **not too difficult** for a computer to calculate (can be done in distributed way)

Questions so far?

Path-based metrics



Which nodes provide the most important **connectivity** or **reachability**?

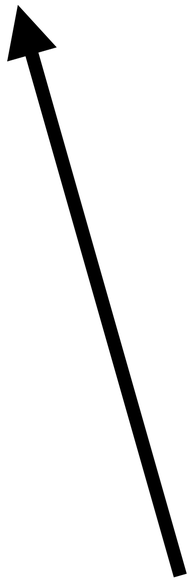
Which nodes, if **removed**, would **damage** the network most?

Less about influence/popularity

Closeness centrality

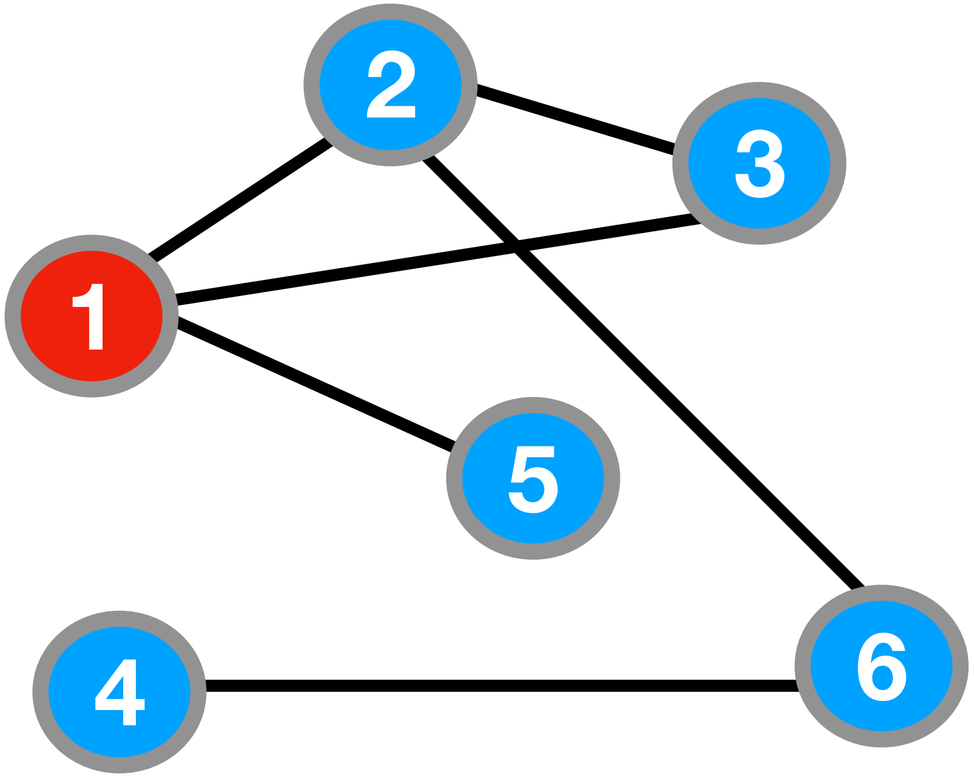
Which node is closest to everything else?

$$C_c(i) = \left[\sum_{j=1}^{N} d(i,j) \right]^{-1}$$



The **smaller** the distances, the **larger** the centrality value

Closeness centrality: example



j	d(1,j)
2	1
3	1
4	3
5	1
6	2

$$C_c(1) = [1 + 1 + 3 + 1 + 2]^{-1} = \frac{1}{8}$$

Closeness centrality: properties

- Usually has **small range** because of short path lengths
- Can be **unstable** — adding or removing a link can dramatically change who is closest
- Path lengths are **expensive** to compute

Betweenness centrality

Which node(s) are most vital for maintaining connectivity?

Number of s**hortest path** routes from **j to k** that go through **i**

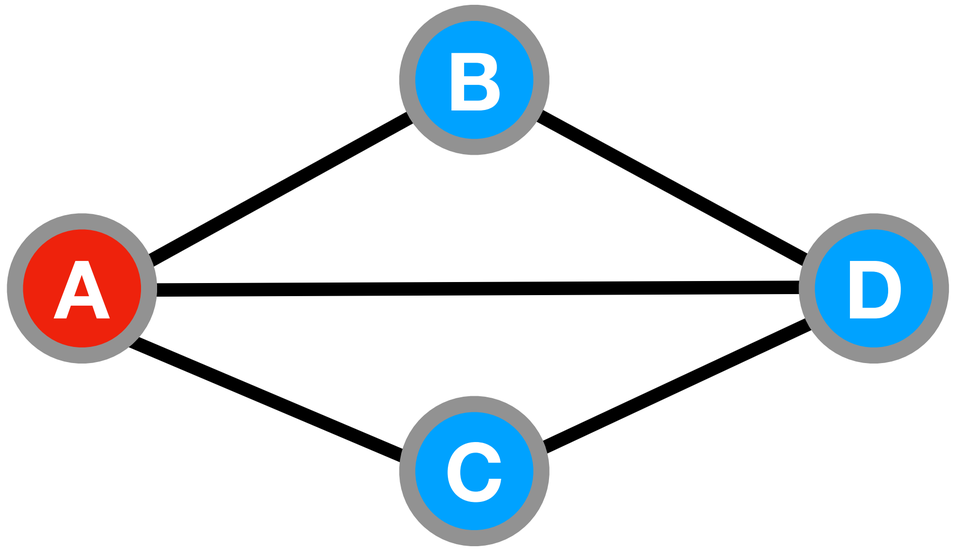
$$C_B(i) = \sum_{j \neq i \neq k} g_{jk}(i) / g_{jk}$$

Sum of all paths apart from those starting or ending at i

Total number of shortest paths from j to k

Betweenness: Example

$$C_B(i) = \sum_{j \neq i \neq k} g_{jk}(i) / g_{jk}$$



Betweenness centrality of A?

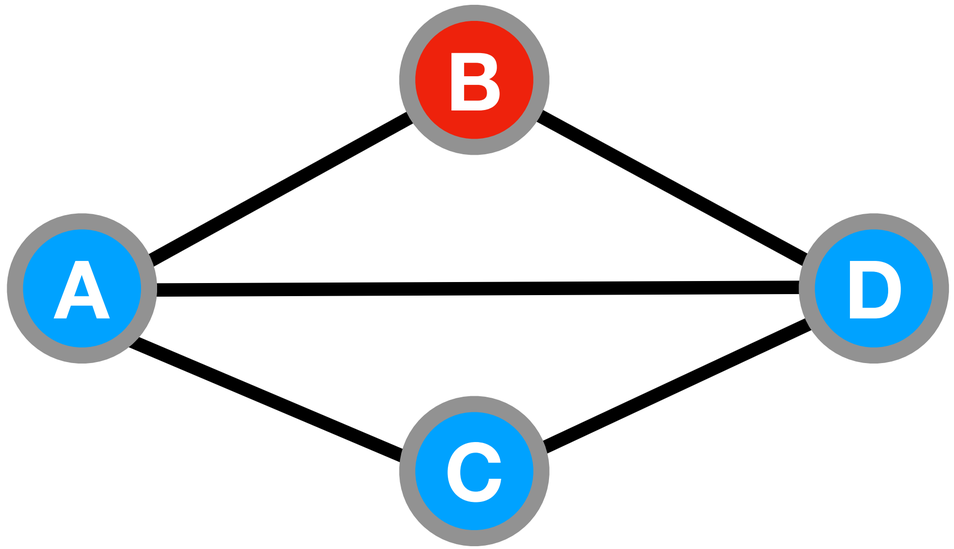
$g_{BC}(A) = 1$ (B-> A -> C)

$g_{BC} = 2$ (B-> A -> C) (B-> D -> C)

$g_{DC}(A) = 0,\ g_{BD}(A) = 0$

$C_B(A) = 1/2$

Betweenness: Example



$$C_B(i) = \sum_{j \neq i \neq k} g_{jk}(i) / g_{jk}$$

Betweenness centrality of B?

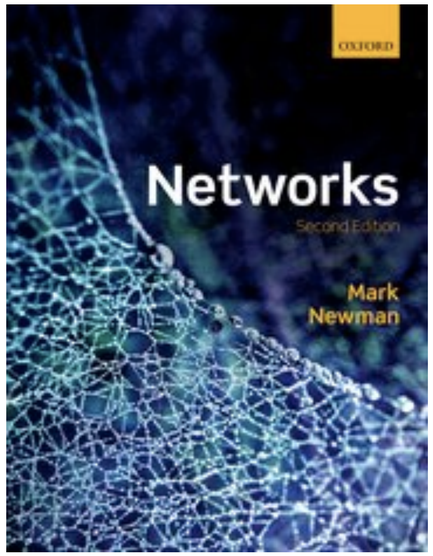
No shortest paths between A, C and D that go through B, so B has betweenness 0.

Betweenness centrality: properties

- Can help identify **vital nodes** for maintaining connectivity
- Used in **Internet monitoring** where (care most about packets getting to destination)
- Finding shortest paths is **expensive** to compute

The more detailed the network, the more complex centrality measures you can devise!

Detailed e.g. directed, weighted, …



See **Networks: an Introduction** by **Mark Newman** for a deeper dive into different centrality measures

FIFA Data Analysis: Jupyter Notebook

[A public data set of spatio-temporal match events in soccer competitions, Luca Pappalardo et al 2018, Nature]

https://www.nature.com/articles/s41597-019-0247-7#Sec9

Eigenvector Centrality

	A	B	C
R1 Nb Sum	3	4	1
R1 Normalise	3/4	1	1/4
R2 Nb Sum	3	1.5	1
R2 Normalise	1	0.5	1/3
R3 Nb Sum	1.5	2	0.5
R3 Normalise	0.75	1	0.25